

DEFINITION

- Global **name table** (= **spelling table** = **identifier table** = **lexeme table**) used to convert scanned identifier names into numeric references.
- Token stores reference to number in addition to or instead of string.
- ASTIdentifier structure will only be interested in number.
- Purpose:
 - Easily identify references to same identifier.
 - Minimize string comparisons to once throughout entire process.

COST

- Cost matters for interpreters
- Cost calculation:
Assume table of n records, m enquiries:
i.e. there are n names in program and they are used m times altogether
i.e. n additions, m enquiries $m > n$
 - S = average cost of 1 search
 - A = cost of 1 addition
 - $\text{Cost} = n A + n S + m S$
(Each addition involves 1 search for duplicates.)

POSSIBLE ORGANISATIONS**Linear List in Chronological Order (default)**

- Keep an table (array) of records, with pointer to last
 - Add to the end
 - Search from the end to beginning
- Cost
 - $S = n/2$ records when record is found
 - $S = n$ when record is not found (for additions)
 - $A = \text{constant } C$
 - $\text{Cost} = nC + n^2 + m.n/2 = O(n(n+m)) = O(n.m)$

Self-Organising List

- In addition to linear list, provide linked list of table indices which moves last used index to front of list
- Costs are same order, but
 - Additional moving costs (constant with linked lists)
 - Real programs cluster usage of identifiers => real search savings

Binary Search Tree in Alphabetical order

- In addition to linear list, provide BST of table indices organizing names in alphabetical order
- Cost
 - $S = \log n$
 - $A = \log n$
 - $\text{Cost} = O((n+m) \log n) = O(m \log n)$
In practise, useful if $n > 50$

Hash Table

- In addition to linear list, provide hash table to organize indices (hash on string)
- Cost: Assume $k =$ size of hash table
 - $S = O(n/k)$
 - $A = O(n/k)$
 - $\text{Cost} = O((n+m)n/k) = O(m.n/k)$
- To reduce Cost, make k big (around $O(m)=100$ is good)